

CS 205M: Lecture 4

Comparing Infinities

and

Hopefully  
(Intro to Propositional Logic)



# Comparing Infinities:

## Countable vs Uncountable Sets

- More precisely, we see how do we compare cardinalities ("sizes") of infinite sets.
- Infinite Sets- Sets containing infinitely many distinct elements.  
E.g.  $\mathbb{N}$ ,  $\{x : x|2\}$  (set of even nos.).
- Comparing cardinalities of finite set: Trivial.  
To check:  $|S| \leq |S'|$ -  
(1) Count no. of elements in  $S$  (say  $k$ ) and,  
(2)  $S'$  (say  $k'$ ). Is  $k \leq k'$ ?



# Abstracting Notion of Comparing Sizes:

- Beyond the notion of counting.

“Universal” Carnival/Mela



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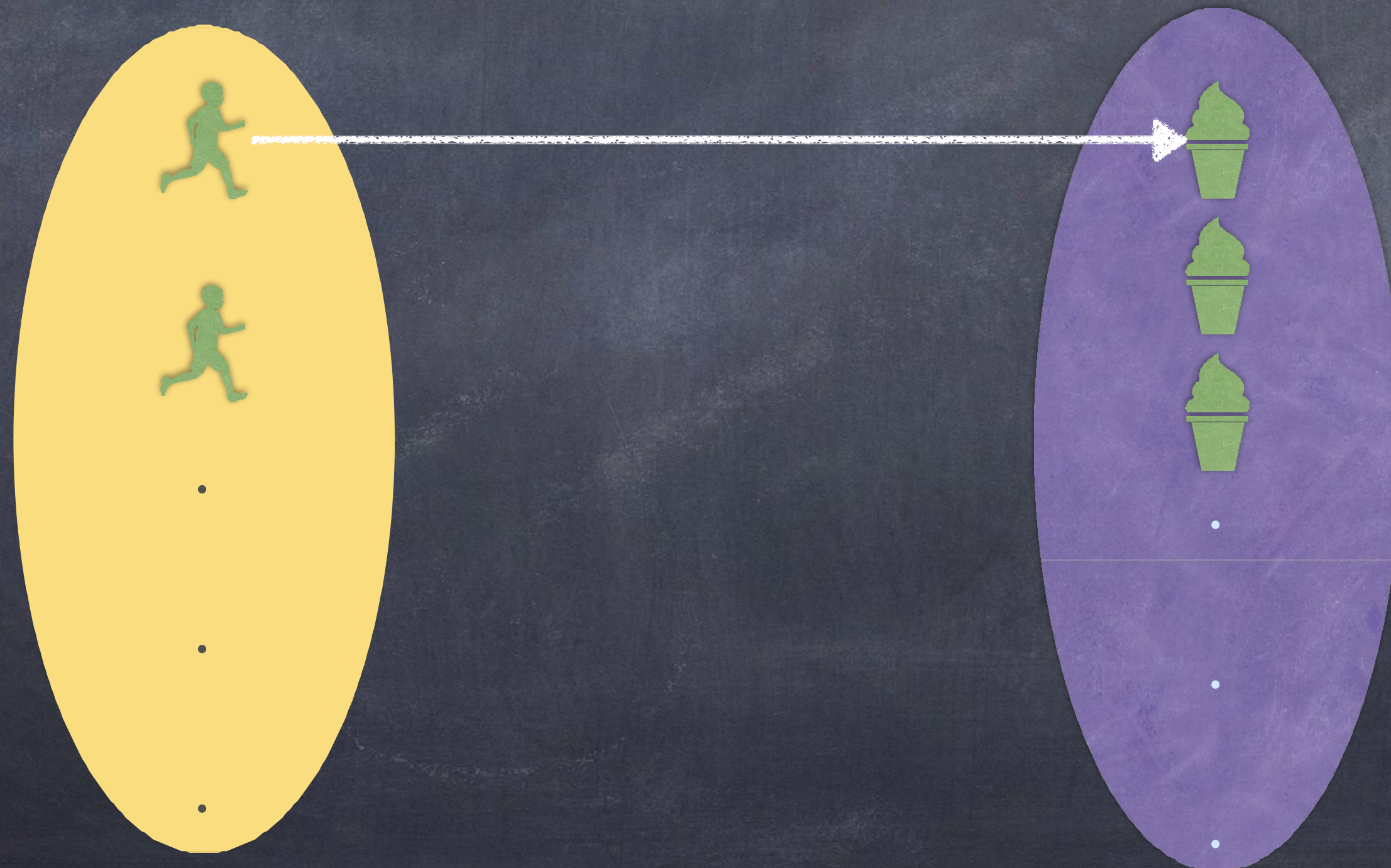




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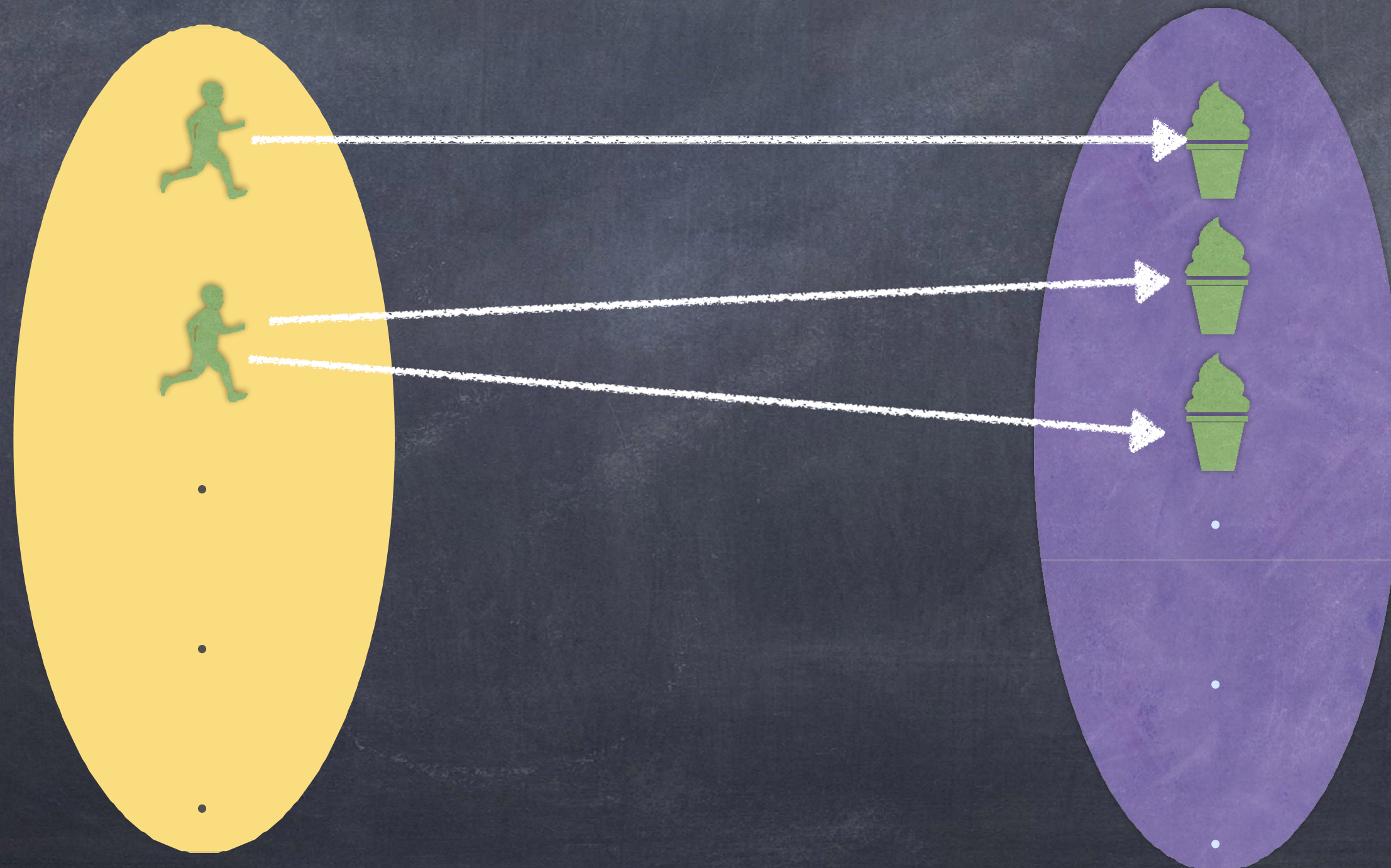




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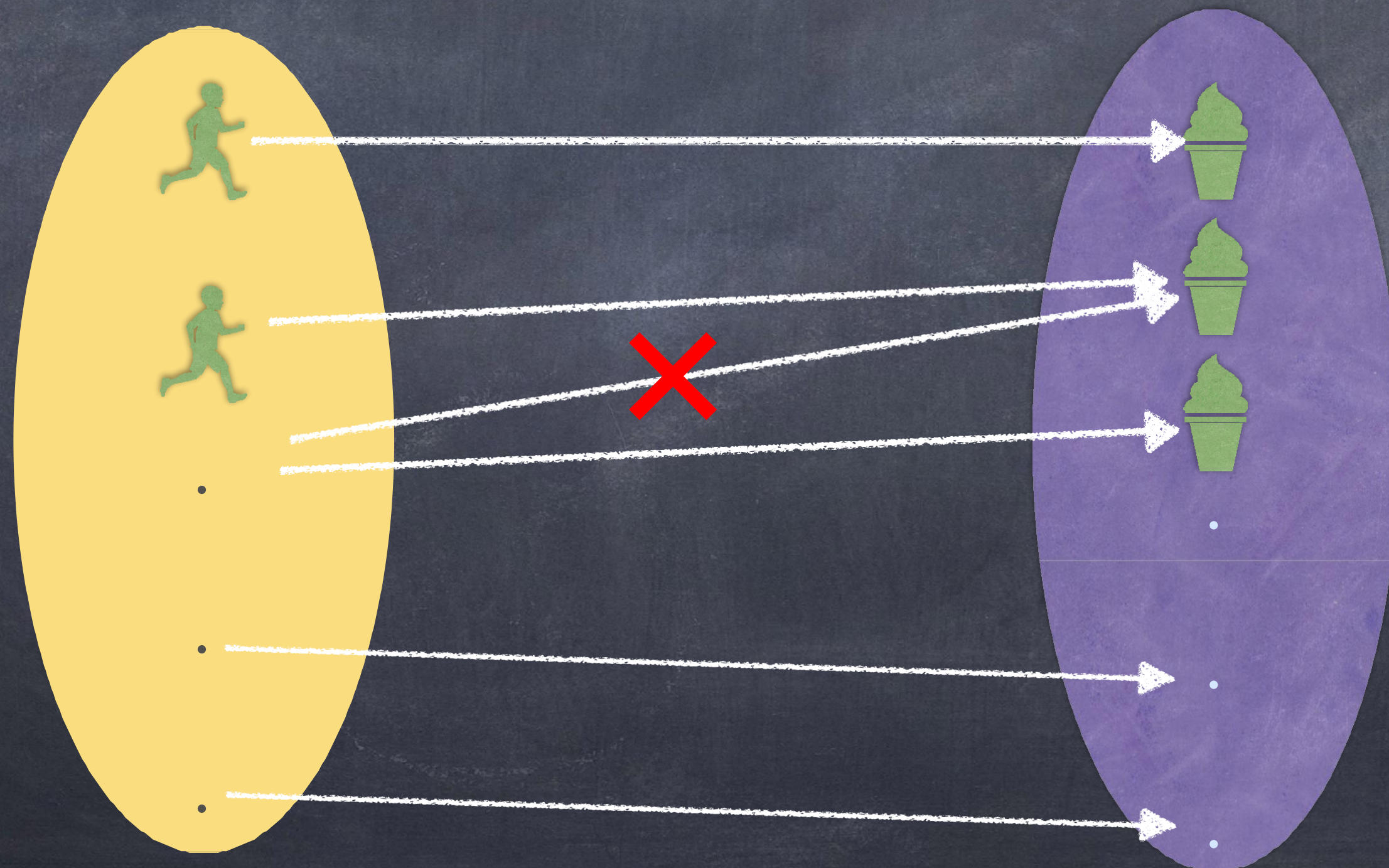




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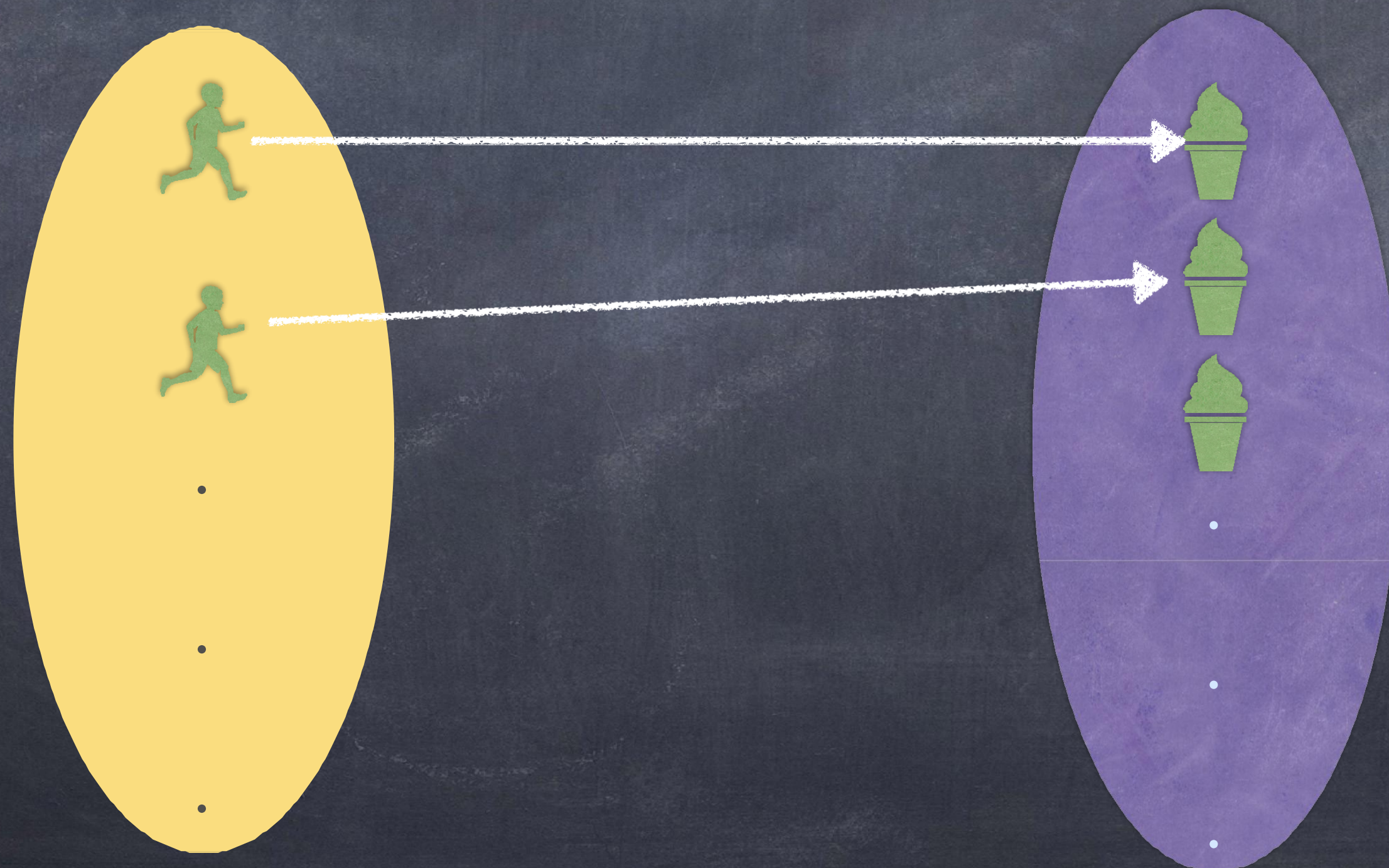


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If there is an injective  
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set of Kids to  
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$$|\text{Kids}| \leq |\text{Ice-creams}|$$





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“Universal” Carnival/Mela

$$| \text{Kids} | \leq | \text{Ice-creams} |$$

$$| \text{Kids} | = | \text{Ice-creams} |$$

iff

$$| \text{Kids} | \leq$$

$$\leq$$

$$| \text{Ice-creams} |$$

and

$$| \text{Ice-creams} | \leq$$

$$\leq$$

$$| \text{Kids} |$$





No. of Naturals vs. No. of Rationals?

$$|\mathbb{N}| < |\mathbb{Q}|?$$



# No. of Naturals vs. No. of Rationals?

- $|\mathbb{N}| \leq |\mathbb{Q}|$  - Trivial - Map every natural no. to itself.
- $|\mathbb{Q}| \leq |\mathbb{N}|$  ?
- $F: \mathbb{Q} \mapsto \mathbb{N}^2$ .  $F(p/q) = (p, q)$ . Hence,  $|\mathbb{Q}| \leq |\mathbb{N}^2|$ .
- $G: \mathbb{N}^2 \mapsto \mathbb{N}$ .  $G((p, q)) = 2^p \times 3^q$ . Notice that if  $(p', q') \neq (p, q)$  then  
$$G((p', q')) \neq G((p, q))$$

(Unique Prime Factorization)
- In fact one could show that for any finite sequence of natural numbers  $(a_1, \dots, a_n)$ , we can associate a unique integer  $2^{a_1} \times 3^{a_2} \times \dots \times k^{a_n}$ , where  $k$  is the  $n^{\text{th}}$  prime.



# No. of Naturals vs. No. of Reals?

$$|\mathbb{N}| \leq |\mathbb{R}|?$$

- $|\mathbb{N}| \leq |\mathbb{R}|$  - Trivial - Map every natural no. to itself.  $n_1$
- $|\mathbb{R}| \leq |\mathbb{N}|$  ? Assume yes. This implies,  $F: [0,1) \mapsto \mathbb{N}$ , and  $F$  is an injective function.

[illegible]



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[illegible]

Add 1 modulo 10.



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[illegible]

Add 1 modulo 10.

The new no. differs from  $i^{th}$  number at  $i^{th}$  digit.

Hence, entirely new number, violating the assumption.



# In General,

- $|X| < |2^X|$ , trivial for finite sets. For infinite sets,
- Let  $X = \{n_1, n_2, \dots\}$

[illegible]

$m_{i,j}$  are either 0 or 1.

# Flip 0s and 1s.



# Why Study Logic?

- **Foundation of mathematics and computer science**
- Helps in:
  - Designing circuits
  - Writing correct software (specifications and verification)
  - Constructing valid arguments and proofs
- Used in AI, algorithms, databases, security protocols, and more



# Propositions.

- **Building Blocks of Logics.** That is why, it is also called as atoms.
- Proposition: A declarative sentence that is either true or false, but not both.
- Examples:
  - "The Earth is round." ✓ (Proposition)
  - " $2 + 2 = 5$ " ✓ (Proposition)
  - "Close the door." ✗ (Not a proposition)
  - " $x + 3 > 5$ " ✗ (Depends on  $x$ ; not a proposition unless  $x$  is defined)



# Logical Operators: $\wedge$ , $\vee$ , $\neg$

Operator	Symbol	Name
NOT	$\neg p$	Negation
AND	$p \wedge q$	Conjunction
OR	$p \vee q$	Disjunction



# Logical Operators: $\wedge$ , $\vee$ , $\neg$

- Examples:
- $\neg(2 < 3) \Rightarrow \text{False}$
- $(2 > 1) \wedge (4 = 4) \Rightarrow \text{True}$
- $(5 > 10) \vee (3 = 3) \Rightarrow \text{True}$



# Logical Operators: $\wedge$ , $\vee$ , $\neg$

- Examples:
- Conjunction (AND) –  $p \wedge q$  : "It is raining and it is cold."  
True only if both are true.
- Disjunction (OR) –  $p \vee q$  : "I will go for a walk or I will stay home."  
True if either (or both) happen.
- Negation (NOT) –  $\neg p$  : "It is not raining."  
True if the statement "it is raining" is false.



# Truth Tables.

p	q	
T	T	
T	F	
F	T	
F	F	



# Truth Tables.

p	q	$\neg p$
T	T	F
T	F	F
F	T	T
F	F	T



# Truth Tables.

p	q	$\neg p$	$p \wedge q$
T	T	F	T
T	F	F	F
F	T	T	F
F	F	T	F



# Truth Tables.

$p$	$q$	$\neg p$	$p \wedge q$	$p \vee q$
T	T	F	T	T
T	F	F	F	T
F	T	T	F	T
F	F	T	F	F



# Implication: $\rightarrow$

- If it rains, then the ground is wet.
- When is the above statement false?



# Implication: $\rightarrow$

- $p \rightarrow q$ : “If  $p$ , then  $q$ ”
- False only when  $p$  is true and  $q$  is false.

$p$	$q$	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T



# Bi-implication, if and only if: $\rightarrow$

- It rains if and only if the ground is wet.
- When is the above statement false?



# Bi-implication, $\leftrightarrow$

- $p \leftrightarrow q$ : “p if and only if q”,  $p \rightarrow q$  and  $q \rightarrow p$
- True when both p and q are same.

p	q	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T



# Propositions to Bits – Bitwise Logic

- Propositions can be represented as bits:  $T \rightarrow 1, F \rightarrow 0$
- Logical operations correspond to bitwise operators.

p	q	$p \wedge q$	$p \vee q$	$\neg p$
1	0	0	1	0
1	1	1	1	0

- Bitwise Ops: AND:  $0101 \& 1100 = 0100$ , OR:  $0101 | 1100 = 1101$ ,  
NOT:  $\sim 0101 = 1010$  (bitwise complement)



# Equivalence of Boolean Expressions:

## Next Class